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## A method of assigning objects to processing units

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### Description

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#### Field of the invention

10 The present invention relates to the field of data processing, and more particularly without limitation, to balancing the assignment of objects in a multi-computing environment.

#### Background and prior art

15 Various multi-computing architectures are known from the prior art where a plurality of processing units is coupled to form a cluster. Such architectures are used in parallel processing and also in the emerging field of blade computing.

Blade computing relies on blade servers, which are modular, single-board computers. An overview of blade computing is given in "Architectures and Infrastructure for Blade Computing", September 2002, Sun microsystems and "THE NEXT WAVE: BLADE SERVER COMPUTING", Sun Microsystems ([www.sun.com/servers/entry/blade](http://www.sun.com/servers/entry/blade)).

20 A content load balancing blade is commercially available from Sun microsystems ("Sun Fire TM B10n). This blade provides traffic and content management functionalities. Content load balancing is achieved based on URLs, CGI scripts and cookies; server

- 5 load balancing is achieved based on server loads, response times, and weighted round-robin algorithms.

US patent application no. 20030105903 shows a web edge server, which comprises a number of blade servers. A switch and an information distribution module are provided for the purpose of balancing. The information distribution module receives an informa-  
10 tion message, performs processing on the message to determine a destination, and forwards a message toward the determined destination via an internal communications network.

### Summary of the invention

The present invention provides for a method of assigning objects to processing units of  
15 a cluster of processing units. Each one of the processing units has a certain storage capacity and load capacity. The storage capacity of a processing unit determines the maximum aggregated size of objects that can be stored by the processing unit. The load capacity of the processing unit determines the maximum processing load that the processing unit can handle.

20 For example, the load capacity of a processing unit can be indicative of the maximum number of access operations the processing unit can provide. Access operations may comprise both read accesses (select) and write accesses (update, insert, delete) to objects stored on the processing unit. For example the load capacity can be expressed as the maximum number of access operations per time unit the processing units can han-  
25 dle.

In order to make maximum usage of the available data processing capacity provided by the cluster of processing units the distribution of the objects over the processing units needs to be balanced. This is accomplished by calculating an index for each object based on object size and object load. For example, the object load is expressed as the  
30 expected mean number of access operations per time unit to the object. The objects are sorted by index in order to provide a sequence.

In the following it assumed without restriction of generality that the sequence is sorted in descending order. In this instance the procedure for assigning of objects to processing units starts with the first object of the sequence. One or more of the objects of the se-  
35 quence are assigned to one processing unit in sequential order until a remaining stor-

5 age capacity and/or a remaining load capacity of that processing unit is too small for consecutive objects of the sequence. When this condition is fulfilled, the procedure is carried out for the next processing unit, whereby the objects that have been previously assigned to the preceding processing unit are deleted from the sequence. This way a minimum number of processing units that are required for handling a given set of ob-  
10 jects can be determined.

In accordance with a preferred embodiment of the invention each processing unit is a single-board computer that has a bus interface to a bus system that couples a plurality of the single-board computers. Each of the single-board computers has its private processing and data storage resources. Data processing tasks or sub-tasks of a complex  
15 data processing task are assigned to the single-board computers by a control unit. The control unit can be a separate hardware unit or a software process that runs on one of the single-board computers. An example of such a distributed data processing system is a cluster of blades.

20 In accordance with a preferred embodiment of the invention the remaining storage capacity of a processing unit is determined by the difference between the storage capacity of the unit and the aggregated size of the objects that have been assigned to the processing unit. Likewise the remaining load capacity of a processing unit is determined by the difference between the load capacity of the unit and the aggregated loads of objects  
25 that have been assigned to the processing unit. On the basis of these definitions of the remaining storage capacity and of the remaining load capacity the minimum number of processing units is determined.

In accordance with a further preferred embodiment of the invention the balancing procedure is performed again in order to further improve the quality of the balancing. For  
30 this purpose the largest gap between the aggregated sizes of objects being assigned to one of the processing units and the largest gap between the aggregated loads of objects being assigned to one of the processing units and the load capacity are determined.

The size gap is divided by the minimum number of processing units and the result of the  
35 division is subtracted from the maximum storage capacity to provide a size threshold level. Likewise, the load gap is divided by the number of processing units and the result

5 of the division is subtracted from the load capacity to provide a load threshold level. When the procedure for assigning the objects to the processing units is performed again, the definition of the remaining storage capacity is the difference between the aggregated size of objects being assigned to the processing unit and the size threshold level whereas the definition of the remaining load capacity is the difference between the aggregated load of the objects being assigned to the processing unit and the load threshold level. As a result of the renewed performance of the assignment procedure, the gap can be substantially reduced.

15 In accordance with a further preferred embodiment of the invention the theoretical storage capacity limit is used as a size threshold. This size threshold is obtained by calculating the difference between the total of the storage capacities of the processing units and the total of the sizes of the objects and dividing the difference by the minimum number of processing units. The result of the division is subtracted from the storage capacity which provides the theoretical limit.

20 Likewise the theoretical load capacity limit is used as a load threshold. This load threshold is obtained by calculating the difference between the total of the load capacities of the processing units and the total of the loads of the objects and dividing the difference by the minimum number of processing units. The result of the division is subtracted from the load capacity which provides the theoretical load capacity limit.

25 On this basis the assignment procedure is performed again whereby the remaining storage capacity is defined as the difference between the aggregated size of the objects of the processing unit and the size threshold whereas the remaining load capacity is defined as the difference between the aggregated load of the objects of the processing units and the load threshold. Typically it will not be possible to assign all of the objects to the minimum number of processing units on this basis. If this is the case one or more iterations are performed.

35 For one iteration an excess amount of memory is divided by the minimum number of processing units. The result of the division is added to the size threshold. Likewise an excess load is divided by the minimum number of processing units. The result of the division is added to the load threshold. On the basis of the incremented size threshold and/or load threshold the assignment procedure is performed again. This process contin-

ues until all objects have been assigned to the minimum number of processing units. This way the quality of the balancing is further improved.

In according with a further preferred embodiment of the invention the size threshold for performing the assignment procedure is varied between the theoretical storage capacity limit and the actual storage capacity. Likewise the load threshold is varied between the theoretical load capacity limit and the actual load capacity. Preferably a new assignment procedure is performed for each permutation of the size threshold/load threshold that can be thus obtained. For each of the resulting assignments of objects to processing units a statistical measure is calculated. This statistical measure is a basis to select one of the assignments for optimal balancing.

In according with a further preferred embodiment of the invention the standard deviation or variance of the sum of the indices of objects assigned to a processing unit is used as a statistical measure. The assignment having the lowest overall quality measure is selected.

In accordance with a preferred embodiment of the invention the object sizes and object loads are normalised for the calculation of the indices. Preferably an index of an object is calculated on the basis of the sum of the normalised object size and normalised object load and the absolute difference of the normalised object size and normalised object load. Preferably the index is obtained by calculating a linear combination of the sum of the normalised object size and normalised object load and the absolute value of the difference of the normalised object size and normalised object load.

In accordance with a preferred embodiment of the invention each one of the processing units is a blade or a blade server. One of the blades can have a program that implements the principles of the present invention in order to perform balancing. This way the number of swap-operations between the blades can be minimised.

In accordance with a preferred embodiment of the invention the principles of the invention are implemented in an application program running on a personal computer. The application program is provided with a list of objects and the estimated sizes and loads of the objects that need to be handled by the cluster of processing units. On the basis of the object sizes and the object loads the minimum number of processing units that are required for the processing task are determined. This information can form the basis for a corresponding investment decision of a customer.

- 5 It is to be noted that the present invention is not restricted to a particular type of object. For example, data objects such as tables, arrays, lists and trees are distributed to processing units, e.g. blades, in accordance with the principles of the present invention. For example, each one of the processing units runs a data processing task to which the respective objects are assigned.

## 10 **Brief description of the drawings**

In the following preferred embodiments of the invention will be described in greater detail by way of example only, by making reference to the drawings in which :

- Figure 1 is a schematic block diagram of a modular computer system having a cluster of blades,
- 15 Figure 2 is illustrative of a flow diagram for assigning of objects to blades and for determining the minimum number of blades,
- Figure 3 is an example for tables that need to be assigned to blades,
- Figure 4 shows the normalised size distribution of the tables of figure 3,
- Figure 5 shows the loads of the tables in terms of the mean number of read accesses per second,
- 20 Figure 6 shows the normalised loads,
- Figure 7 shows the normalised sizes and loads per table,
- Figure 8 shows a sorted sequence of the tables,
- Figure 9 shows the sorted sequence of the tables illustrating the normalised sizes and loads of each table,
- 25 Figures 10 to 14 show the assignment of tables to blade 1 of the cluster of blades
- Figures 15 and 16 show the assignment of tables to blade 2,
- 30 Figure 17 shows the resulting assignment of tables to blades that is obtained on the basis of the method of figure 2,

- 5 Figure 18 is illustrative of a preferred embodiment of the invention where the procedure of figure 2 is performed again with lower size and load thresholds,
- Figure 19 is illustrative of a preferred embodiment of a method of the invention where the load and size thresholds are varied iteratively,
- 10 Figure 20 is illustrative of a further preferred embodiment of the invention where the size and load thresholds are varied in predetermined steps,
- Figure 21 is illustrative of a computer system performing the assignment of objects to blades.

### Detailed description

- 15 Figure 1 shows cluster 100 of blades B1, B2, B3, ... BN. Each one of the blades has processor 102 and memory 104. In the example considered here, all memories 104 have the same storage capacity and all blades have the same load capacity. The blades are coupled by a network 106, such as a bus system. The number N of blades of cluster 100 needs to be chosen such that a given number of M objects of varying sizes and loads can be handled.
- 20 For example, cluster 100 implements a so-called search engine. In this instance identical search processes run on each one of the blades. The assignment of data objects, such as index tables, to blades can be stored in a dispatcher unit (not shown on the drawing) of cluster 100. This way data objects are assigned to blades and data processing tasks running on the blades.
- 25 Figure 2 shows a procedure for assigning the objects to blades and to thereby determine the minimum value for N.
- In step 200 an sorting index is calculated for each one of the M objects. An sorting index of an object is indicative of the amount of blade resources the object requires. The sorting index serves to sort the objects in decreasing order of blade resource requirements.
- 30 For example the sorting index is calculated on the basis of the sum of the normalised object load and normalised object size plus the absolute value of the difference of the normalised load and size or a linear combination thereof.

5 In step 201 a sorting operation is performed in order to sort the M objects by sorting index. The corresponding object sequence is provided in step 202. In step 204 the index i for the blades is initialised to 1.

In step 206 processing of the object sequence starts with the first object of the sequence, i.e. the object having the largest sorting index value. The first object of the sequence is assigned to a first one of the blades, i.e. blade B1, in step 206. In step 208  
10 the first object that has been assigned to blade B1 is deleted from the sequence.

In step 210 the sizes of the objects that have already been assigned to blade B1 are summed up in order to provide an aggregated object size of blades B1. Next the size of a gap  $G_S$  between the aggregated object size of blade B1 and a size threshold  $T_S$  is  
15 calculated. When the assignment procedure of figure 2 is carried out the first time, the size threshold  $T_S$  is the storage capacity of one of the blades.

In step 211 the loads of the objects that have already been assigned to blade B1 are summed up in order to provide an aggregated load of blade B1. Next a gap  $G_L$  between the aggregated object loads of blade B1 and a load threshold  $T_L$  is calculated. When the  
20 assignment procedure of figure 2 is carried out the first time, the load threshold  $T_L$  is the load capacity of one of the blades.

In step 212 it is determined whether there is a next object in the ordered sequence that fits into both gaps  $G_S$  and  $G_L$ . In other words, a consecutive object following the first object in the object sequence that has an object size small enough to fit into gap  $G_S$  and  
25 at the same time has an object load that is small enough to fit into  $G_L$  is searched.

The next consecutive object in the sequence that fulfils this condition is assigned to blade B1 in step 214 and deleted from the sequence in step 216 before the control goes back to step 210.

If there is no such object that fulfils the condition of step 212, step 218 is carried out. In  
30 step 218 it is determined whether all objects have already been assigned to blades. In other words, in step 218 it is checked whether the sequence is empty. If this is not the case the index i is incremented in step 220 and the control goes back to step 206 in order to assign remaining objects of the sequence to the next blade B2.

If the contrary is the case the index i is the minimum number N of blades that are re-  
35 quired to handle the M objects, i.e.  $i = N$ . This number is output in step 220. The mini-



5     mum number N of blades that are required to handle the M objects can be a basis for an investment decision for purchasing of a corresponding number of blades. The assignment of objects to blades is output in step 224 in order to visualise the quality of the object size balancing.

Figure 3 shows an example of objects to be assigned to blades. In the example considered here the objects are a number of 20 different tables having various sizes between  
10     50MB and 3,566MB as indicated in figure 3. For example, table 1 has a size of 3,250MB, table 2 has 250MB, table 3 has 750MB, etc. The table sizes can be actual table sizes or average table sizes that have been obtained by monitoring a real-life data processing system. Alternatively the table sizes are estimates for the purpose of planning cluster 100 (cf. figure 1).  
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Figure 4 shows the normalised table sizes of the tables of figure 3. The normalised size of a table is obtained by dividing the table size by the maximum storage capacity of one of the blades.

Figure 5 shows the loads of the tables. In the preferred embodiment considered here  
20     the mean number of access operations (selects) per second that are performed or expected to be performed with respect to a table is used as a load indicator. In the example considered here the load as expressed by the number of read accesses per second to a table varies between five read accesses per second (table 6) and 950 read accesses per second (table 8). The table loads given in figure 5 can be actual table loads  
25     or average table loads that have been obtained by monitoring a real-life data processing system. Alternatively the table loads are estimates or simulation results for the purpose of planning cluster 100.

Figure 6 shows the normalised table loads. The normalised table load of a given table is obtained by dividing the number of read accesses per second of that table by the load  
30     capacity of one of the blades.

In the preferred embodiment considered here the storage capacity of a blade is 4,096MB. Hence a normalised table size of one indicates that the table has the absolute maximum size that can be handled by a given blade hardware.

The load capacity of a blade is the maximum possible access load that can be handled  
35     by a core engine running on one of the blades in the example considered here. This

5 maximum value can be determined by benchmarks, by experiment or simulation. The load capacity depends on various parameters such as hardware and software characteristics and network bandwidth if a network is used to perform the table accesses. In the preferred embodiment considered here, the load capacity of one of the blades is 1,000 read accesses per second. For the purpose of explanation only read accesses  
10 are considered here. However, other typical data processing tasks, such as accesses that involve changes to the data, can also be taken into consideration for determining load capacity and table loads.

Figure 7 shows a diagram illustrating the various normalised table sizes and table loads. As apparent from figure 7 there are tables that require only a minimum amount of blade  
15 resources, especially table 13, that has a small table size and a low table load. On the other hand there are tables that require a large amount of blade resources as they have both large table size and high table load. The distribution of tables to blades aims to balance the usage of the respective blade resources with respect to both size and load. The two constraints, i.e. size and load, are combined into a common sorting index that  
20 provides an objective measure for the amount of blade resources a given table requires.

The following definition of the sorting index is used for the purposes of explanation only and without restriction of generality:

Sorting index =  $W1 * (size + load) + W2 * \text{absolute value} (size - load)$ ,

where size is the table size,

25 load is the table load

W1 is a weighting factor for (size + load) and

W2 is a weighting factor for the absolute value of the difference of size and load.

For the purposes of the following example the weighting factors W1 and W2 are set to one without restriction of generality.

30 In this case the above expression evaluates as follows:

If size > load: sorting index = 2 \* size

If size = load: sorting index = 2 \* size = 2 \* load

5 If size < load: sorting index = 2 \* load.

Figure 8 shows a sorted sequence in descending order of the sorting indices of the tables. The table that has the largest sorting index value and thus requires the largest amount of blade resources is table 8, whereas the table having the smallest sorting index value and which does require the lowest amount of blade resources is table 13 as  
10 apparent from figure 8.

Figure 9 shows the sorted sequence of figure 8 illustrating the normalised table sizes and table loads of the individual tables.

Figure 10 illustrates the first iteration for assigning tables to blades based on the ordered sequence as shown in figure 8 and 9. It is assumed that each one of the blades  
15 has a storage capacity that corresponds to a normalised size threshold  $T_S = 1$  and a load capacity that corresponds to a normalised load threshold  $T_L$ .

The assignment procedure starts with the first table of the sorted sequence, i.e. table 8. Table 8 is assigned to blade 1 as illustrated in figure 10. This results in a normalised gap  $G_S$  of unused storage capacity  $G_S = 0,6948$  and a gap  $G_L$  of unused load capacity  
20  $G_L = 0,05$ .

Next consecutive tables in the ordered sequence are searched that have table sizes and table loads that fit into the respective gaps  $G_S$  and  $G_L$ . These are tables 6 and 13. As table 6 precedes table 13 in the ordered sequence, it is assigned to blade 1 as illustrated in figure 11.

25 Next the aggregated table size and the aggregated table load of blade 1 is updated as illustrated in figure 12. This results in a updated normalised gap  $G_S = 0,6362$  and an updated normalised gap  $G_L = 0,045$  as illustrated in Fig. 13. Now tables of the ordered sequence that are consecutive to table 8 that fit into both gaps  $G_S$  and  $G_L$  are searched.

The only table that fulfils both conditions is table 13 which is thus assigned to blade 1 as  
30 illustrated in figure 13. The result is shown in figure 14. The updated normalised gap  $G_S$  is 0,6167 and the updated normalised gap  $G_L$  is 0,03. No table within the ordered sequence satisfies both constraints.

As there remain unassigned tables in the sequence an additional blade 2 is required as illustrated in figure 15. By definition the first table of the remaining ordered sequence,

5 i.e. table 20 is assigned to blade 2. This results in normalised gaps  $G_S = 0,1294$  and  $G_L = 0,5600$ . The only table that satisfies both constraints is table 12 which is thus moved onto blade 2 as illustrated in figure 15.

Figure 16 illustrates the result of this assignment. The updated normalised gap  $G_S$  is 0,0024 and the updated normalised gap  $G_L$  is 0,4500. None of the tables of the remain-  
10 ing ordered sequence satisfies these constraints such that an additional blade 3 is required etc. This procedure is carried out until the ordered sequence is empty, i.e. after all tables have been assigned to one blade. The resulting assignment of blades and the resulting minimum number of blades is illustrated in figure 17. In the example considered here the minimum number  $N$  of blades that is required for handling of all of the ta-  
15 bles is  $N = 10$ .

In order to further improve the quality of the balancing the method of figure 18 can be carried out on the basis of the result provided by the method of figure 2. In step 800 the largest remaining gaps,  $G_S$  and  $G_L$  are determined as illustrated in figure 17. In the ex-  
20 ample considered in figure 17 the largest remaining gap  $G_S$  is on blade 1 and the largest remaining gap  $G_L$  is on blade 10.

In step 802 the largest remaining gap  $G_S$  is divided by  $N$  which yields delta 1 and the largest remaining gap  $G_L$  is divided by  $N$  which yields delta 2.

In step 804 the size threshold  $T_S$  is reduced by delta 1 and the load threshold  $T_L$  is reduced by delta 2. In step 806 the method of figure 2 is performed again with the re-  
25 duced thresholds as determined in step 804. In case the sorting indices and the original object sequence obtained from the method of figure 2 by performing steps 200 to 202 have been saved, these steps can be skipped in order to start processing of the method of figure 2 in step 204. The result is a further improved balancing.

Figure 19 shows an alternative approach for refining the balancing. In step 900 delta 3  
30 is calculated by calculating the difference of the sum of the storage capacity of the blades and the sum of the object sizes of all objects to be assigned to the blades and by dividing the difference by the minimum number of blades  $N$ . In step 902 delta 4 is calculated by calculating the difference of the sum of the load capacities of the blades and the sum of the object loads of all objects to be assigned to the blades and by dividing  
35 the difference by the minimum number of blades  $N$ . In the example considered here no normalized values are used in steps 900 and 902.

5 In step 904 the size threshold  $T_S$  is reduced by the normalised value of delta 3. The normalised value of delta 3 is obtained by dividing delta 3 by the storage capacity of one of the blades.

Likewise the load threshold  $T_L$  is updated in step 906 by the normalised delta 4. Normalisation of delta 4 is performed by dividing delta 4 by the load capacity of one of the  
10 blades.

The reduced size threshold  $T_S$  and the reduced load threshold  $T_L$  correspond to the theoretical limit of blade resources that are required for handling of the given objects. As the object granularity is finite the theoretical threshold limits will be surpassed in most cases:

15 In order to refine the balancing the method of figure 2 is performed again (step 908) on the basis of the size threshold  $T_S$  obtained from step 904 and the reduced load threshold  $T_L$  obtained from step 906 with the modification that the number of blades for the renewed performance of the method of figure 2 is limited by the minimum number of blades  $N$  that has been determined previously. In other words, the assignment procedure of figure 2 stops when the minimum number  $N$  of blades has been used up; this  
20 may have the consequence that there will be an excess amount of memory and/or load requirement for the last blade of the minimum number  $N$  of blades considered for the assignment procedure that surpasses the respective thresholds  $T_S$  and/or  $T_L$ . Again steps 200 to 202 of the method of figure 2 can be skipped in case the sorting indices  
25 and the sorted sequence has been stored when the method of figure 2 was performed the first time. In step 910 it is determined whether there is such an excess amount of memory requirement and/or load requirement for the last blade  $N$ , i.e. blade 10. If this is not the case the resulting assignment of objects to blades is output in step 912.

If the contrary is the case step 914 is carried out in order to calculate the values of delta  
30 5 and/or delta 6. Delta 5 is obtained by dividing the excess amount of memory, if any, by the number of blades  $N$ . Likewise delta 6 is obtained by dividing the excess load requirement, if any, by the number of blades  $N$ .

On this basis the size and/load thresholds are incremented in step 916. From there the control goes back to step 908.

- 5 Steps 908 to 916 are carried out repeatedly until there is no longer an excess amount of memory and/or load requirement that cannot be provided by the given number N of blades.

Figure 20 shows a further alternative for refinement of the balancing. First steps 900 and 902 of the method of figure 19 are carried out in order to obtain delta 3 and delta 4.

- 10 The gaps delta3 and delta4 are now scanned by stepwise variation of the respective size and load thresholds in order to identify a more balanced assignment of objects to blades. For this purpose delta 3 and delta 4 are divided into equidistant intervals having step sizes delta 7 and delta 8, respectively, as shown in step 1000. The number of increments for performing the respective scans can be predefined or is user selectable.

- 15 It is to be noted that the number of increments for scanning  $G_S$  and for scanning  $G_L$  does not need to be the same. Preferably the scans are performed independently from each other such that the total number of assignments that is considered is the number of increments for the  $G_S$  scan multiplied by the number of increments for the  $G_L$  scan.

- 20 In step 1002 the size and load thresholds are set to the respective theoretical minima that are required to provide sufficient blade resources for handling of the given number of objects. On this basis the method of figure 2 is performed again in step 1004. In step 1006 a statistical measure is calculated as a quality measure for the balancing of the assignment of objects to blades obtained as a result of step 1006, if any. For example, the standard deviation of the aggregated sorting index values of objects assigned to
- 25 each one of the blades is calculated. In other words, for each blade the total of the sorting index values of the objects that have been assigned to the blade is calculated. This provides one total sorting index value per blade. Next the standard deviation is calculated for the total sorting index values.

- 30 In step 1008 at least one of the thresholds  $T_S$  or  $T_L$  is incremented by the normalised value of delta 7 or the normalised value of delta 8, respectively. Next step 1004 is carried out again on the basis of the incremented size and load thresholds. Steps 1004 to 1008 are carried out repeatedly until the respective scans through  $G_S$  and  $G_L$  have been completed and the corresponding assignments of objects to blades have been obtained. In step 1010 one of the assignments is selected based on the statistical measures. For
- 35 example, the assignment having the lowest standard deviation is selected.

5 It is to be noted that this procedure is limited by the minimum number of blades N. For assignments that do not fit on this given minimum number of blades N no statistical measure needs to be calculated as these assignments are not considered further in the procedure.

10 Figure 21 shows a computer 108 that has processor 110 for running program 112. Program 112 has module 113 for calculating a sorting index value for each object to be assigned to a blade. Further program 112 has module 114 for sorting of the objects by sorting index value and module 116 for assigning of objects to blades.

15 Further computer 108 has storage 118 for storing a table listing the objects, object sizes, and object loads of objects to be assigned to blades, storage 120 for storage of a storage capacity value of the blades, storage 121 for storage of a load capacity value of the blades, and storage 122 for storing of the number of blades. Further computer 108 has interface 124 for coupling to workstation 126.

20 In operation the table with the object names/numbers, object sizes and object loads is entered via interface 124 and stored in storage 118. Further a storage capacity value for the storage capacity of each individual blade is entered via interface 124 and stored in storage 120. Likewise a load capacity value for the load capacity of each individual blade is entered via interface 124 and stored in storage 121.

25 Next program 112 is invoked. Program 112 calculates a sorting index for each object contained in the table stored in storage 118 on the basis of the normalised object size and object load. The resulting indices are entered into the table stored in storage 118 by module 113. Next module 114 sorts the table of storage 118 by decreasing storage index to provide a sequence of objects (cf. the sequence of figure 8). Next module 116 performs the method of figure 2 in order to determine the minimum number of required blades.

30 This minimum number is stored in storage 122 and is output via user interface 124. This number can be a basis for a users investment decision for purchasing the number of blades to realise the data processing system being capable of handling the objects as listed in the table.

35 In addition, module 116 can perform the methods of figure 18, 19 and/or 20 for refining of the balancing.

- 5 Alternatively, computer 108 is one of the blades. In this instance computer 108 can dynamically change the assignment of objects to blades when the object size changes. This way frequent swapping operations can be avoided. In other words, the creation of "hot spots" is prevented by balancing the load.



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## List of Reference Numerals

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	100	cluster
	102	processor
10	104	memory
	106	network
	108	computer
	110	processor
	112	Program
15	113	module
	114	module
	116	module
	118	storage
	120	storage
20	121	storage
	122	storage
	124	interface
	126	workstation